



AFRL

Applying Categorical Thinking to Practical Domains

Jared Culbertson

Air Force Research Laboratory - Autonomous Capabilities Team (ACT3)

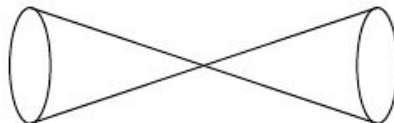
(Joint work with Kirk Sturtz, Dan Guralnik, Jakob Hansen, Dan Koditschek, Paul Gustafson, Peter Stillér
and influenced by conversations with many other researchers)

Background: Algebraic geometry & Representation theory

G reductive complex algebraic group with Lie algebra \mathfrak{g}

$$\mathcal{N} = \{x \in \mathfrak{g} \mid x \text{ is nilpotent in every representation of } \mathfrak{g}\}$$

Example: $\mathfrak{g} = \mathfrak{sl}_2$; $\mathcal{N} = \{2 \times 2 \text{ matrices with zero trace and determinant}\}$



Proved theorems like...

Theorem (C. 2010)

Given functors $\mathcal{D}' \xrightarrow{i_*} \mathcal{D} \xrightarrow{j^*} \mathcal{D}''$ and t -structures on \mathcal{D}' and \mathcal{D}'' such that

- (G1) j^* is essentially surjective
- (G2) there is a t -structure on $\widetilde{i_* \mathcal{D}'}$ induced by the one on \mathcal{D}'
- (G3) the t -structures on \mathcal{D}' and \mathcal{D}'' are compatible
- (G4) $j^* f = 0 \implies f$ factors through $\widetilde{i_* \mathcal{D}'}$
- (G5) $\widetilde{i_* \mathcal{D}'} = \{X \mid j^* X = 0\}$

then there is a unique t -structure on \mathcal{D} extending those on \mathcal{D}' and \mathcal{D}''



Theorem (C. 2010)

For each orbit $C \subset \mathcal{N}$, there is a fully faithful functor

$$\mathcal{IC}(C, -) : \mathcal{M}_{\text{Poi}G}(C) \rightarrow \mathcal{M}_{\text{Poi}G}(\mathcal{N})$$

and every simple perverse Poisson sheaf is of the form $\mathcal{IC}(C, \mathcal{F})$

Theorem (C. 2010)

For a Poisson sheaf $\mathcal{F} = \mathcal{O}_C \otimes \mathcal{L}$ on an orbit C , we have

$$i_{C'}^{! \text{Poi}} H^i(\mathcal{IC}(C, \mathcal{F})) = \mathcal{O}_{C'} \otimes \mathcal{L}''$$

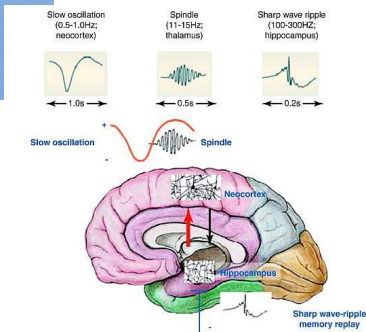
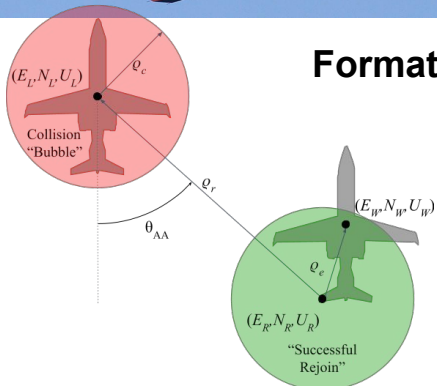
AFRL - Autonomous Capabilities Team



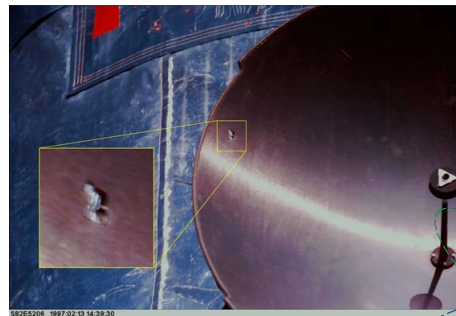
X-62A Modified F-16 (VISTA)



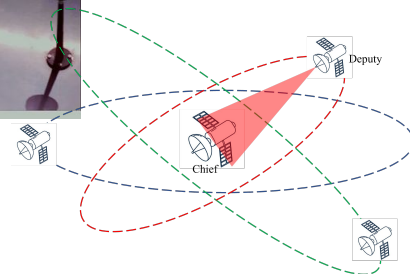
Formation Flying



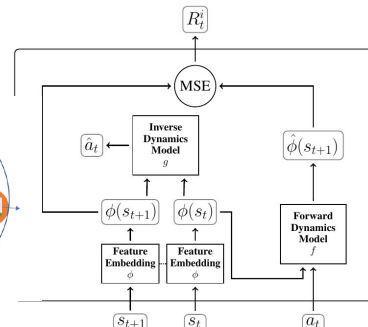
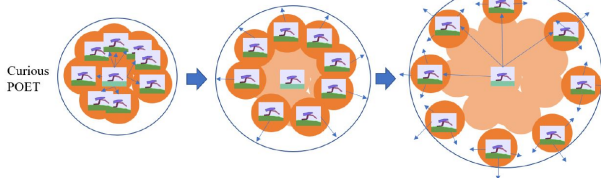
Neuro-mechanisms and architectures



Images: NASA



Multi-agent Spacecraft Inspection



Intrinsic Curiosity & Open-ended Learning

Categorical Probabilities (1960s+ : Lawvere, Huber, Giry, Doberkat, ...)

The Probability Monad $\mathcal{P}: \mathit{Meas} \rightarrow \mathit{Meas}$

$$\begin{array}{l} \text{:ob} \quad X \quad \mapsto \quad \mathcal{P}X = \{\text{all probability measures on } X\} \\ \text{:ar} \quad X \xrightarrow{f} Y \quad \mapsto \quad \underbrace{\mathcal{P}X \xrightarrow{\mathcal{P}f} \mathcal{P}Y}_{P \mapsto P f^{-1}} \end{array}$$

Kleisli Category of \mathcal{P}

$\mathit{Meas}_{\mathcal{G}}$

- Objects: measurable spaces
- Morphisms: stochastic kernels $X \times \Sigma_Y \rightarrow [0, 1]$
- Composition: given $X \xrightarrow{f} Y \xrightarrow{g} Z$

$$(g \circ f)(x, C) = \int_Y g_C df_x$$

Many results of the flavor: “ * can be expressed as...”

Probability Theory	$Meas_{\mathcal{G}}$
Probability measure P	Morphism $1 \xrightarrow{P} X$
Measurable function $X \xrightarrow{f} Y$	$x \mapsto \delta_{f(x)} = \text{Dirac measure}$
U -valued random variable	Composition $1 \xrightarrow{P} \Omega \xrightarrow{\delta_X} U$
Regular conditional probability	$Meas_{\mathcal{G}}$ morphism
Conditional expectation	Diagram in $Meas_{\mathcal{G}}$
Stochastic process	Functor $\mathcal{T} \rightarrow Meas_{\mathcal{G}}$

Theorem: (C, Sturtz, 2013)
 Bayesian probability can be realized in $Meas_{\mathcal{G}}$

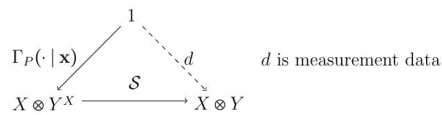
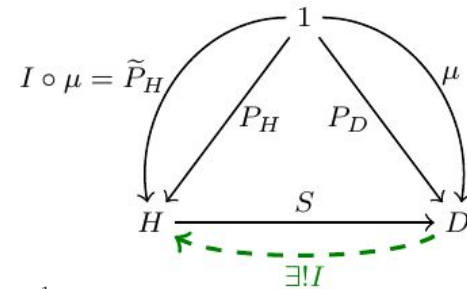


Figure 17: The generic nonparametric Bayesian model for stochastic processes.

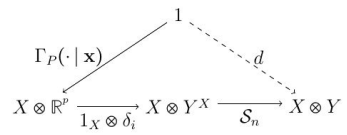
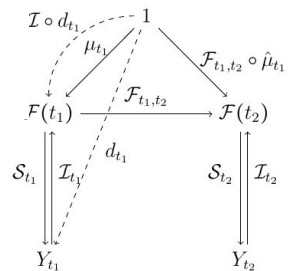


Figure 25: The generic parametric Bayesian model.



[C., Sturtz, A Categorical Foundation for Bayesian Probability, ACS, 2014]

Figure 40: The hidden Markov model viewed in \mathcal{P} .

Original work

- Extremely useful for me in translating to a language I understood
- Not particularly immediately impactful on algorithmic development

Unsatisfying conclusions, but recently...

A free energy principle for generic quantum systems

Chris Fields^{a*}, Karl Friston^b, James F. Glazebrook^{c, d} and Michael Levin^e

^a 23 Rue des Lavandières, 11160 Caunes Minervois, FRANCE

^b Wellcome Centre for Human Neuroimaging, University College London, London, WC1N 3AR, UK

^c Department of Mathematics and Computer Science, Eastern Illinois University, Charleston, IL 61920 USA

^d Adjunct Faculty, Department of Mathematics, University of Illinois at Urbana-Champaign, Urbana, IL 61801 USA

^e Allen Discovery Center at Tufts University, Medford, MA 02155 USA

January 3, 2022

Category Theory in Machine Learning

Dan Shiebler

University of Oxford

daniel.shiebler@kellog.ox.ac.uk

Bruno Gavranović

University of Strathclyde

bruno@brunogavranovic.com

Paul Wilson

University of Southampton

paul@statusfailed.com

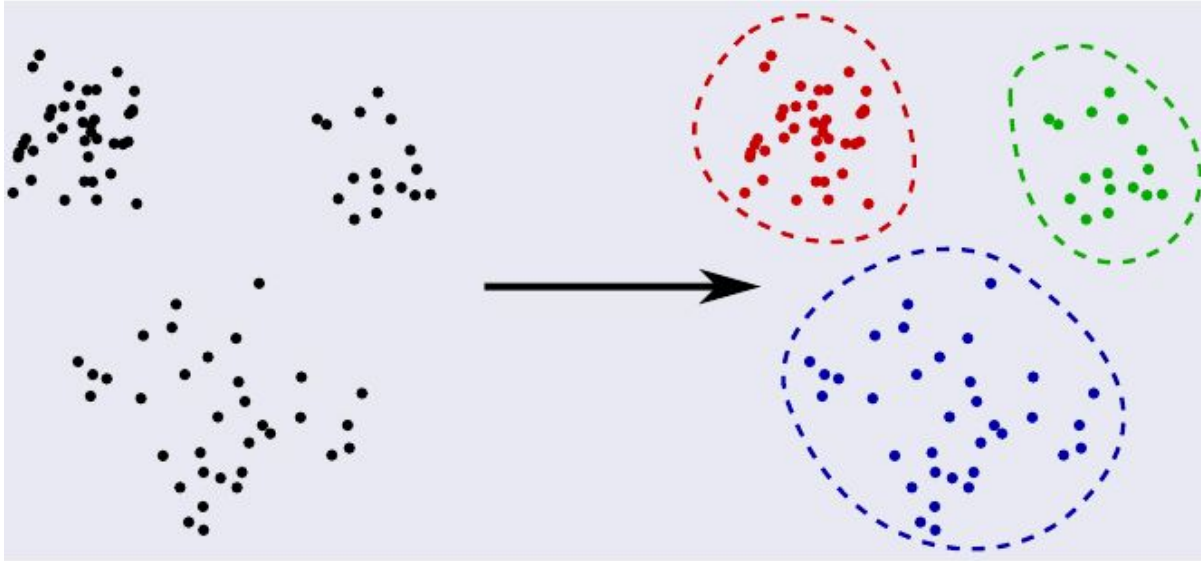
Conditional Distributions for Quantum Systems

Arthur J. Parzygnat

Institut des Hautes Études Scientifiques
Bures-sur-Yvette, France

parzygnat@ihes.fr

Data Clustering

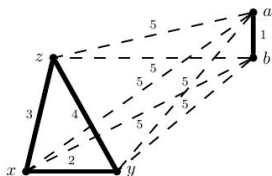


Kleinberg (2002): There are no *consistent* methods of clustering with partitions

- ┌───┐
└───┘
- Rich (all partitions realizable)
 - Scale-invariant
 - Intra/inter-cluster distance transformation invariant

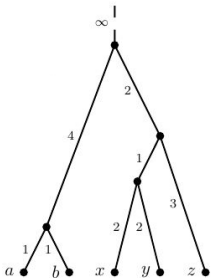
Hierarchical Clustering

Category of Metric Spaces



SL

Category of Ultrametric Spaces



Morphisms: non-expansive maps

$$d(f(x), f(y)) \leq d(x, y)$$

Carlsson–Memoli (2008): There is a unique *functorial* method (single-linkage clustering) for assigning a hierarchical clustering to a metric space.

Q: Are there other useful categories of metric spaces with meaningful projections?

Theorem (C., Guralnik, Stiller, 2016):

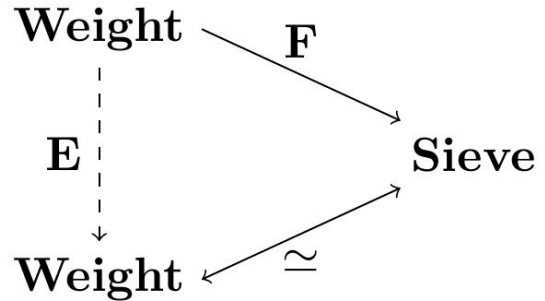
There is an equivalence of categories $\text{Weight} \simeq \text{Sieve}$.

Fibers are directed complete posets under ordering induced from identity maps

{ “nice” subcategories of **weight spaces** }

Composition with functor to Weight satisfies an idempotency condition

{ “stationary” structuring maps to **sieves** }



Corollary: *There is no functorial clustering projection to cut metrics*

Reality: almost no one uses this theory in practice

Did not address:

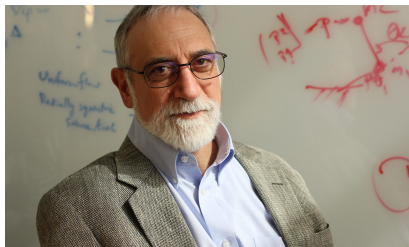
- Cluster uncertainties
- Outlier detection/removal
- **Algorithm efficiency**
- Choice of underlying metrics
- **Software implementations**
- Convincing examples of value
- Characterizations of useful clustering domains

Mismatch with practitioners' needs

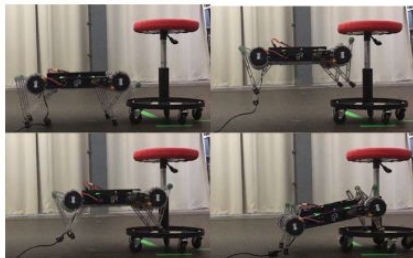
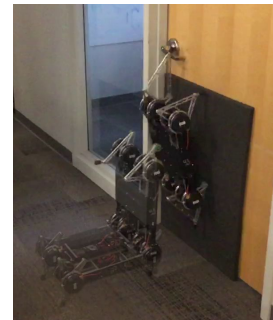
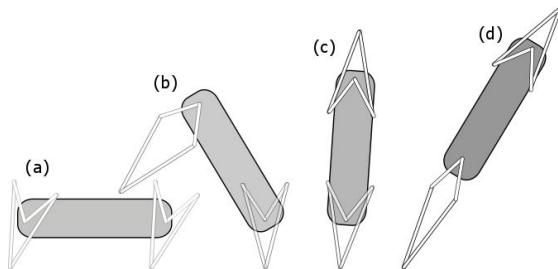
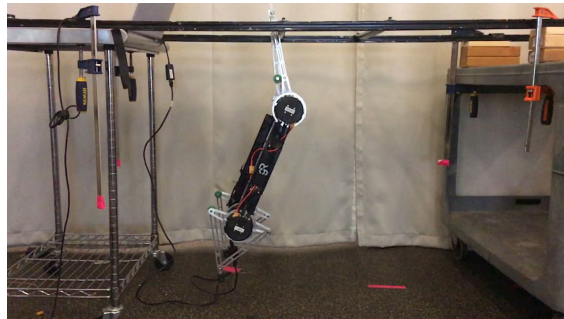
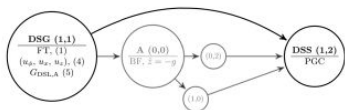
- Opaque categorical language with no clear interface for users
- **Consistency in this sense is rarely considered in applications**
- Existing tools are often sufficient for exploratory analysis
- Simplifications were guided by mathematical considerations

Lesson: “Nice theory → Applications” rarely brings near-term value

Compositional problems in robotics



 **Penn**
UNIVERSITY OF PENNSYLVANIA **Dan Koditschek**



[R. Burridge, et al, "Sequential Composition of Dynamically Dexterous Robot Behaviors," IJRR 1999.]

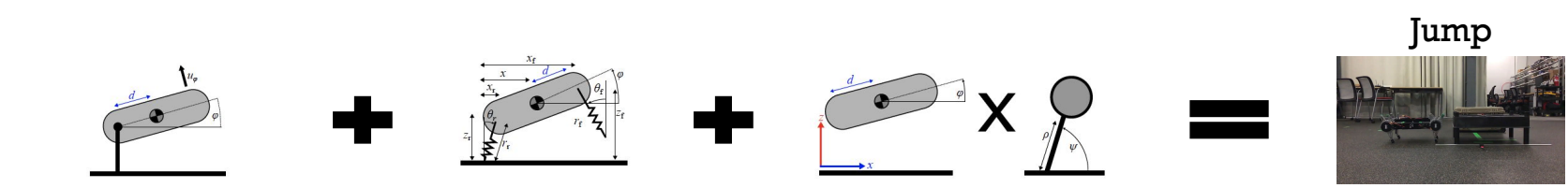
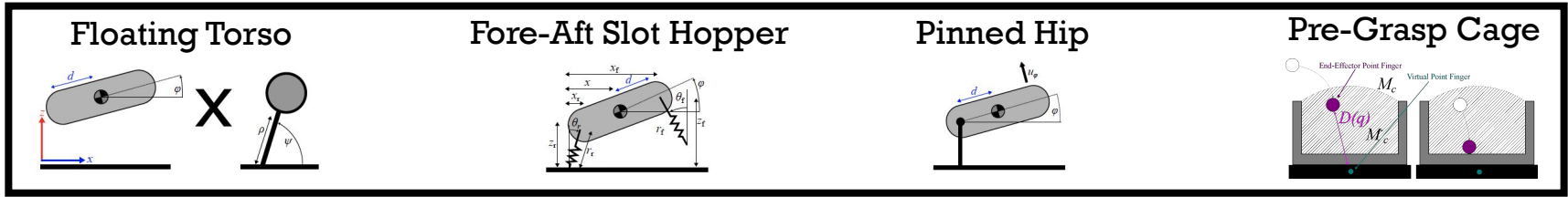
[A. De, et al, "Parallel Composition of Templates for Tail-Energized Hopping" ICRA 2015.]

[A. De, et al, "A Universal Template for Pitch Steady Behaviors in Planar Floating-Torso Locomotion Models," In Prep, 2020.]

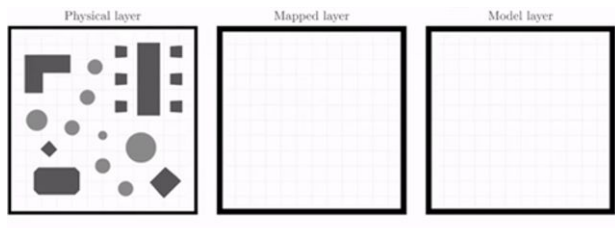
[V. Vasilopoulos, et al, "Sensor-Based Reactive Execution of Symbolic Rearrangement Plans by a Legged Mobile Manipulator" IROS 2018.]

[T. Topping, et al, "Quasi-Static and Dynamic Mismatch for Door Opening and Stair Climbing" ICRA 2017.]

Emerging Calculus of Behaviors



Transformations

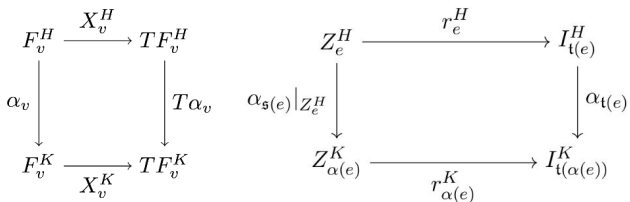


[Vasilopoulos, Koditschek, 2018]



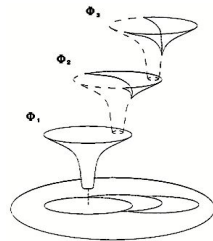
A **hybrid semiconjugacy** $\alpha: H \rightarrow K$ is:

- (1) a graph morphism $\alpha: G(H) \rightarrow G(K)$
- (2) maps of active sets $\alpha_v: I_v^H \rightarrow I_v^K$



Hybrid semiconjugacies

Sequential Composition



[Burridge, Rizzi, Koditschek, 1999]



A **directed system** $H_i \xrightarrow{H} H_f$ consists of

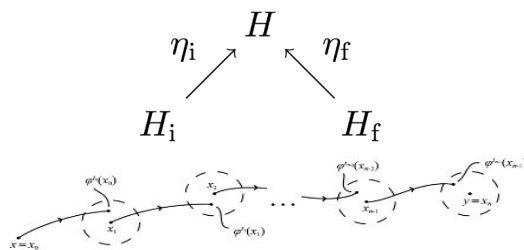
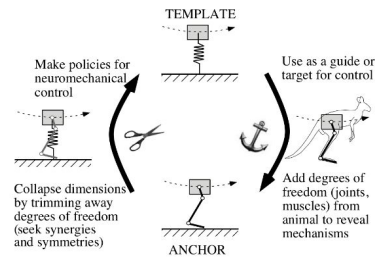


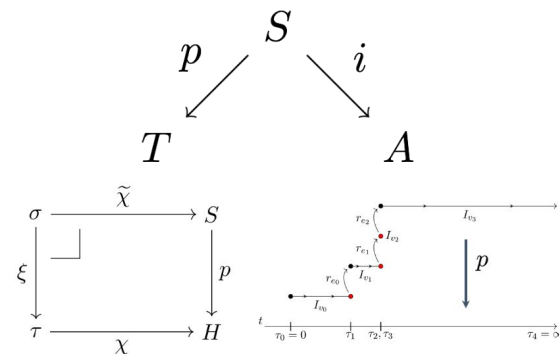
Image source: Alongi and Nelson, *Recurrence and Topology*. AMS, 2007.

Directed systems

Templates & Anchors



[Koditschek, Full, 1999]



Hybrid subdivisions

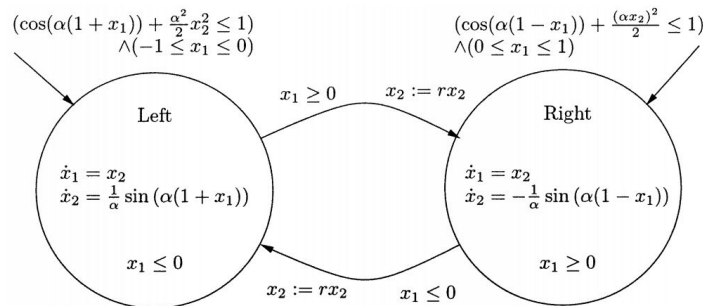
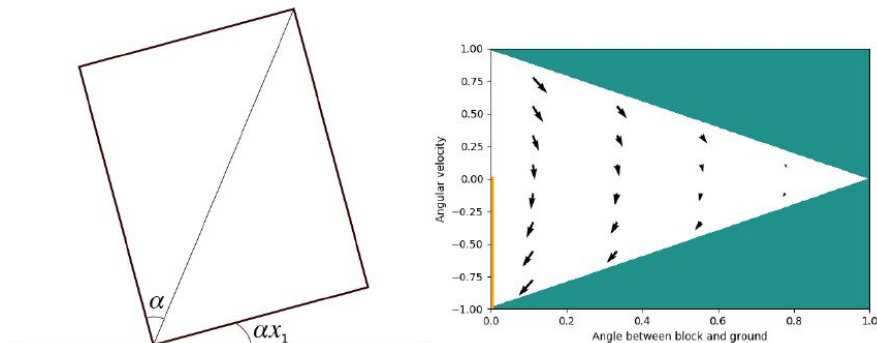
A hybrid system H consists of

1. a directed graph* $G = (V, E, \mathfrak{s}, \mathfrak{t})$;

2. for each continuous mode $v \in V$,

- an ambient manifold M_v
- a vector field X_v on M_v
- an active set $I_v \subset M_v$
- a flow set $F_v \subset I_v$

3. for each reset $e \in E$, a guard set $Z_e \subset I_{\mathfrak{s}(e)}$ and an associated reset map $r_e: Z_e \rightarrow I_{\mathfrak{t}(e)}$.



[Image: Lygeros et al., "Dynamical properties of hybrid automata.," 2003].

* We work here in (a cat. equiv. to) the category of directed reflexive graphs.

A hybrid semiconjugacy $\alpha: H \rightarrow K$ is:

(1) a graph morphism $\alpha: G(H) \rightarrow G(K)$

(2) maps of active sets $\alpha_v: I_v^H \rightarrow I_v^K$

$$\begin{array}{ccc}
 F_v^H & \xrightarrow{X_v^H} & TF_v^H \\
 \alpha_v \downarrow & & \downarrow T\alpha_v \\
 F_v^K & \xrightarrow{X_v^K} & TF_v^K
 \end{array}$$

Flow Compatibility Condition

$$\begin{array}{ccc}
 Z_e^H & \xrightarrow{r_e^H} & I_{t(e)}^H \\
 \alpha_{s(e)}|_{Z_e^H} \downarrow & & \downarrow \alpha_{t(e)} \\
 Z_{\alpha(e)}^K & \xrightarrow{r_{\alpha(e)}^K} & I_{t(\alpha(e))}^K
 \end{array}$$

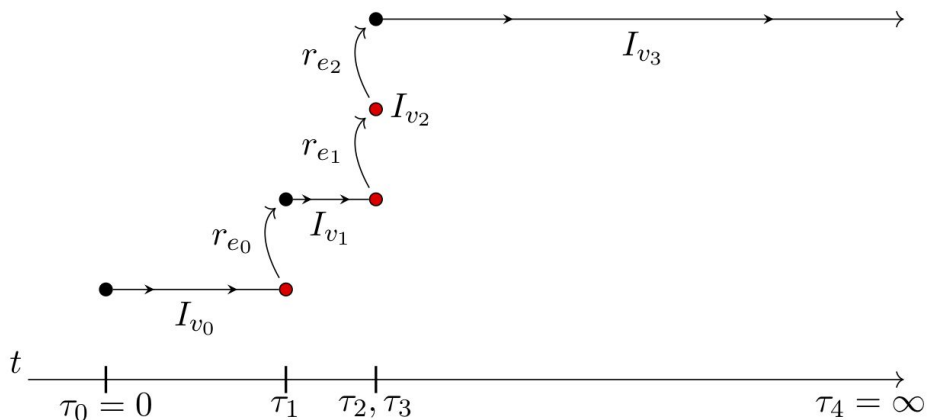
Reset Compatibility Condition

Notes:

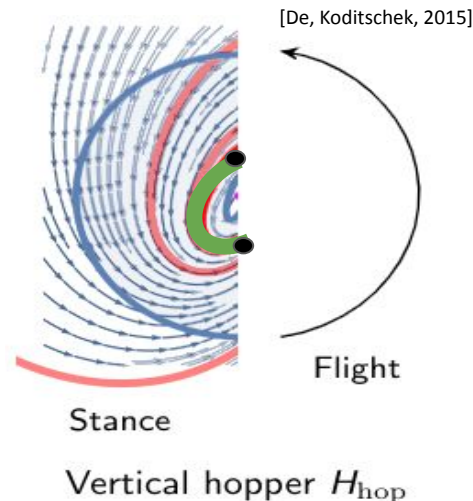
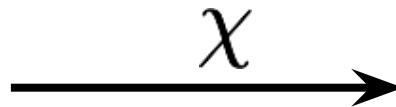
- Conjugacy, submersions, embeddings straightforward
- Technical tools necessary to allow for practically-relevant subsets of manifolds (smooth sets)

$$\text{Hom}_{\mathbf{S}}(A, B) = \text{colim}_{U \supseteq A} \{f \in \text{Hom}_{\mathbf{M}}(U, M_A) \mid f(A) \subset B\}$$

Hybrid executions as semiconjugacies



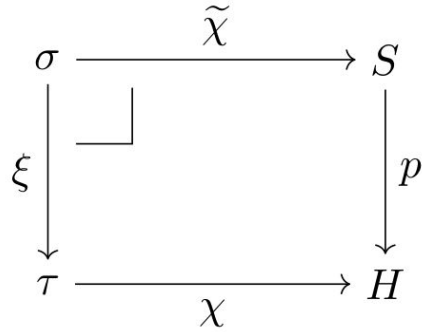
Hybrid Time Trajectory



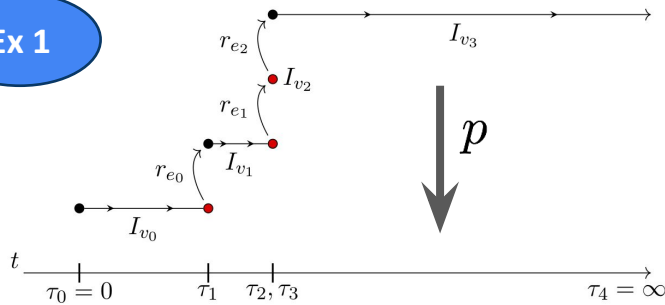
Theorem (Pappas, Lerman, CGSK)

Semiconjugacies from hybrid time trajectories encode all executions of the codomain system.

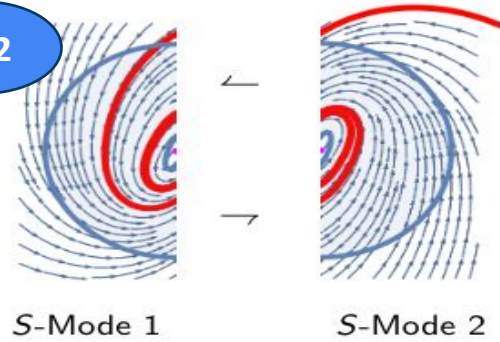
A **subdivision** is a hybrid submersion where



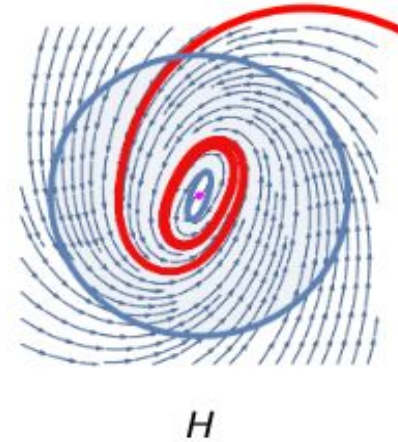
Ex 1



Ex 2



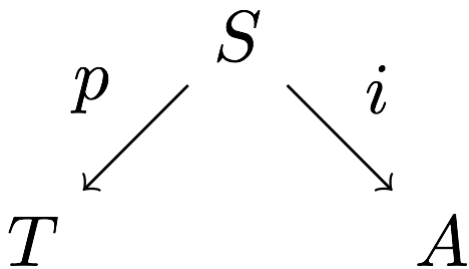
[De, Koditschek, 2015]



Slicing a Continuous Systems

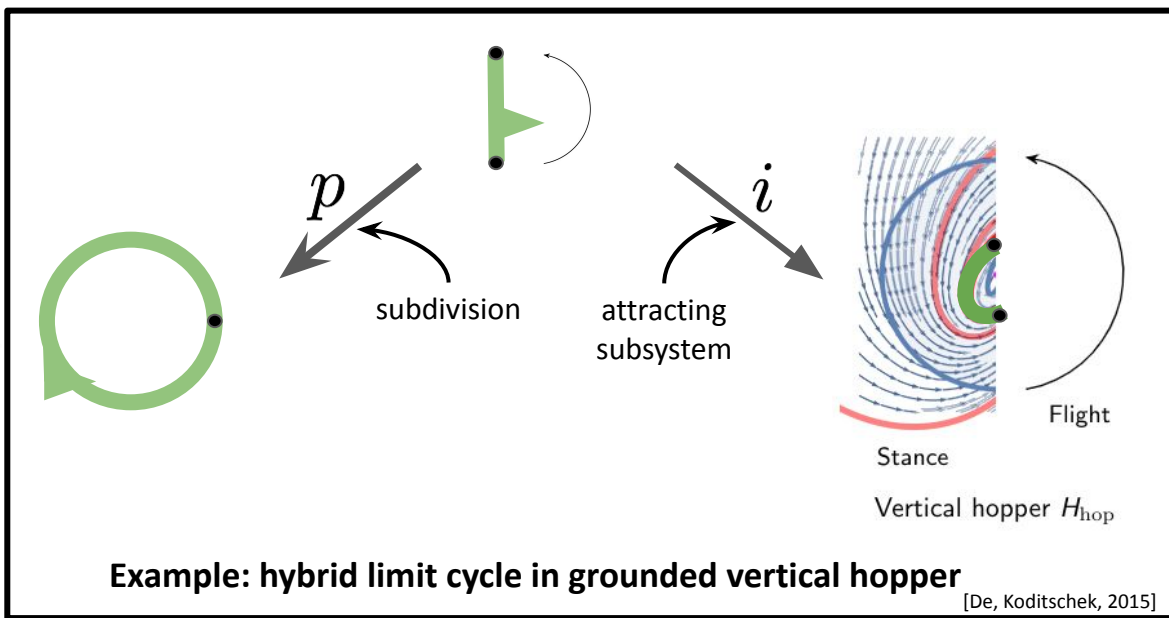
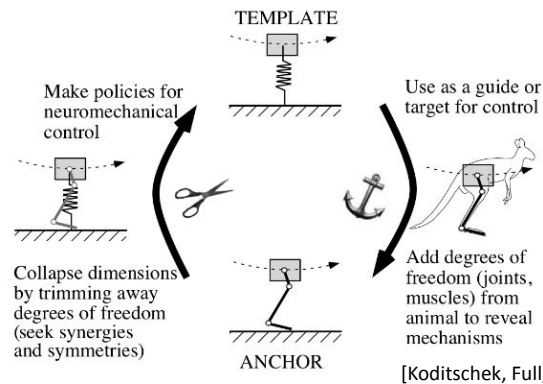
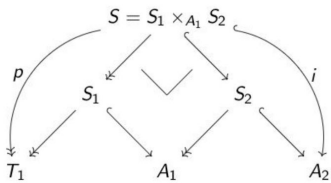
A **template-anchor pair** is a span $T \xleftarrow{p} S \xrightarrow{i} A$ such that

- ▶ p is a hybrid subdivision;
- ▶ i is a hybrid embedding;
- ▶ $i(S)$ is attracting in A .



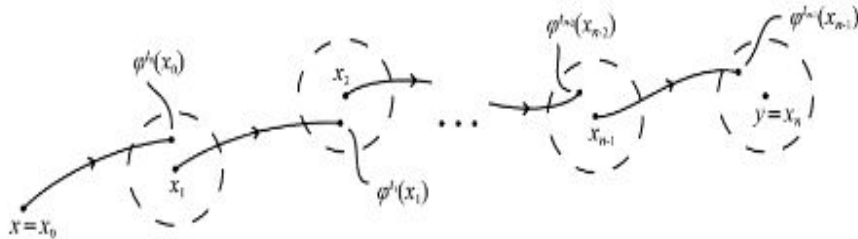
Template-anchor relationship as a span

Theorem (CGKS). Template-anchor pairs are weakly associatively composable.



Conley's "Fundamental Theorem"

Theorem (Norton): *Every flow on a compact metric space decomposes into a chain-recurrent part and a gradient-like part.*



Classical (ϵ, T) -chains

Image source: Alongi and Nelson, *Recurrence and Topology*. AMS, 2007.

Theorem (Gustafson, Kvalheim, Koditschek, 2019)

Every hybrid system with compact state space satisfying a guard-set contraction property decomposes into a chain-recurrent part and a gradient-like part.

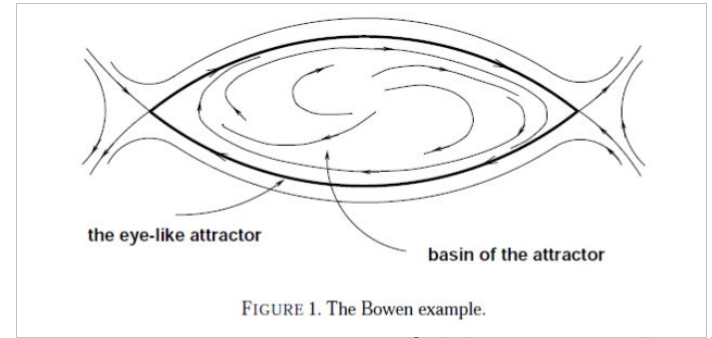


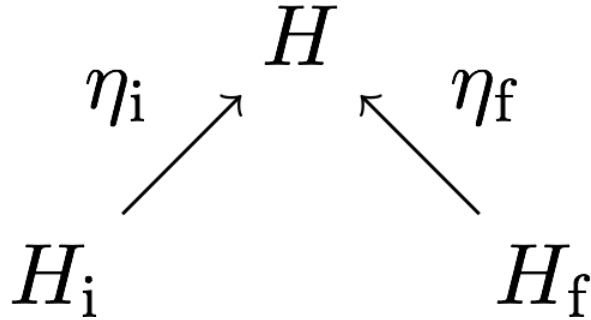
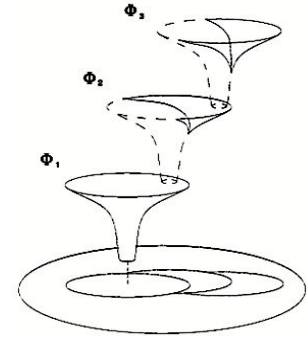
FIGURE 1. The Bowen example.

[Baladi, Bonatti, Bernard, 1999]

A **hybrid (ϵ, T) -chain** is roughly an execution with allowable ϵ -jumps

- after each reset
- after flowing for $> T$ time

A **directed** system $H_i \xrightarrow{H} H_f$ consists of



Initial subsystem

Final subsystem

where H_f is an (ϵ, T) -sink in H

Key observations

1. State spaces not sufficient interfaces
 - Initial and final systems include dynamics
2. Standard executions do not compose well
 - (ϵ, T) -chains do compose!
 - Generalizes topological or measure theoretic approaches
3. "Same" system could have many applications
 - Different initial or final subsystems
 - Specifying new interface into system

Structural Results \longrightarrow Double Category of Hybrid Systems

Sequential composition

$$\begin{array}{ccc}
 \begin{array}{ccc} M & \xrightarrow{H_1} & N \\ \downarrow f & \zeta_1 \Downarrow & \downarrow g \\ P & \xrightarrow{H'_1} & Q \end{array} & \odot & \begin{array}{ccc} N & \xrightarrow{H_2} & K \\ \downarrow g & \zeta_2 \Downarrow & \downarrow h \\ Q & \xrightarrow{H'_2} & L \end{array} \\
 & & = \begin{array}{ccc} M & \xrightarrow{H_2 \odot H_1} & K \\ \downarrow f & \zeta_2 \odot \zeta_1 \Downarrow & \downarrow h \\ P & \xrightarrow{H'_2 \odot H'_1} & L \end{array}
 \end{array}$$

$$\begin{array}{ccc}
 P & \xrightarrow{H'_1} & Q \\ \downarrow f' & \zeta'_1 \Downarrow & \downarrow g' \\ S & \xrightarrow{H''_1} & T
 \end{array}$$

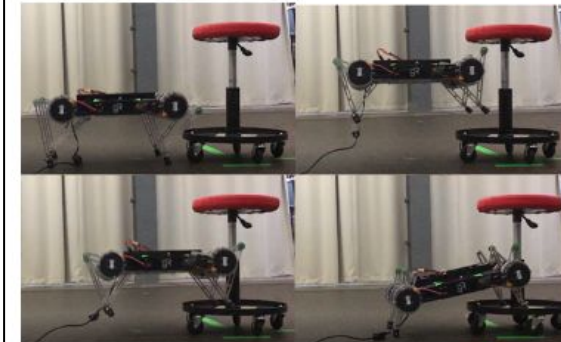
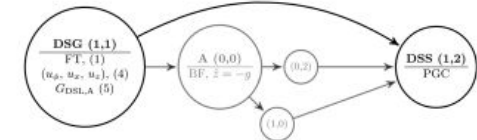
$$\begin{array}{ccc}
 M & \xrightarrow{H_1} & N \\ \downarrow f' \circ f & \zeta'_1 \circ \zeta_1 \Downarrow & \downarrow g' \circ g \\ S & \xrightarrow{H''_1} & T
 \end{array}$$

Category H:

- ▶ objects are hybrid systems
- ▶ vertical category encodes hybrid semiconjugacy
- ▶ horizontal category encodes sequential composition
- ▶ has products and fiber products along hybrid submersions
- ▶ fibered over category of graphs
- ▶ framed bicategory if we allow reset relations
- ▶ associative template composition via spans

Theorem (CGKS)

Hybrid systems, directed systems and hybrid semiconjugacies form a cartesian and cocartesian double category .



[Example: Topping, Vasilopoulos, De, Koditschek, 2019]

Hybrid semiconjugacy composition

Simple Type Theory for Mobile Manipulation

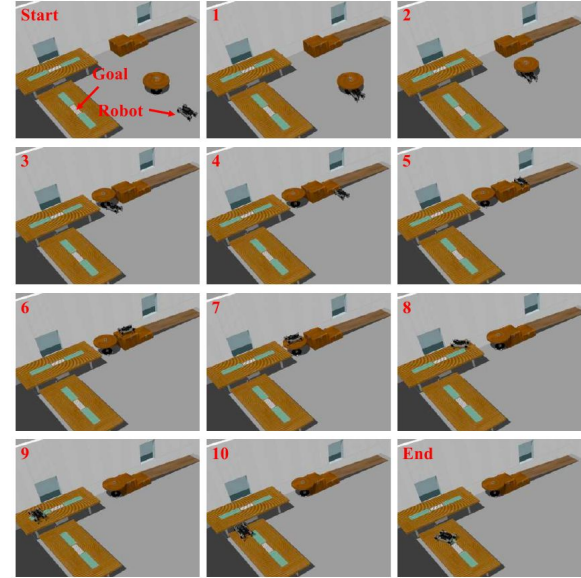
Type	Steady-state behavior	
Free	no navigational goals	“Free type theory” on robot behaviors
Object	stay near the object	
Target _x	maintain position at $x \in X$	
Nest	inhabit the nest	

Operational Semantics



Type	Continuous system (M, X)	Morphism	Reactive hybrid system
Free	$(\mathbb{R}^2, 0)$	$A \rightarrow \text{Object}$	$F_{\text{nav}}(A) \rightarrow (\mathbb{R}^2, X_d) \rightarrow (B_\epsilon(\Delta), X_d _{B_\epsilon(\Delta)})$
Object	$(B_\epsilon(\Delta), X_d _{B_\epsilon(\Delta)})$	$A \rightarrow \text{Target}_x$	$F_{\text{nav}}(A) \rightarrow (\mathbb{R}^2, X_t) \rightarrow (B_\epsilon(x), X_t _{B_\epsilon(x)})$
Target _x	$(B_\epsilon(x), X_t _{B_\epsilon(x)})$	$A \rightarrow \text{Nest}$	$F_{\text{nav}}(A) \rightarrow (\mathbb{R}^2, X_n) \rightarrow (B_\epsilon(r^*), X_n _{B_\epsilon(r^*)})$
Nest	$(B_\epsilon(r^*), X_n _{B_\epsilon(r^*)})$	$A \rightarrow \text{Free}$	$F_{\text{nav}}(A) \rightarrow (\mathbb{R}^2, 0)$

[Gustafson, C., Koditschek, Hybrid dynamical type theories for navigation, 2021]



[Image:Vasilopoulos, 2021]

Implementation: <https://github.com/PaulGustafson/HybSys>

Choosing the right abstractions and simplifications

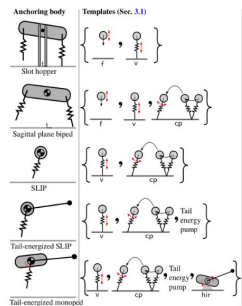
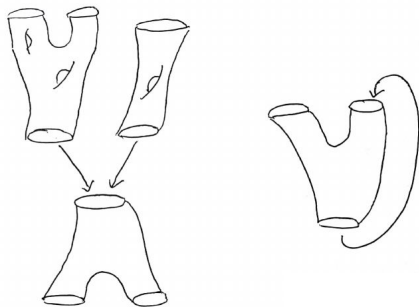
Examples:

1. More abstract and elegant formulations meet grounded approaches

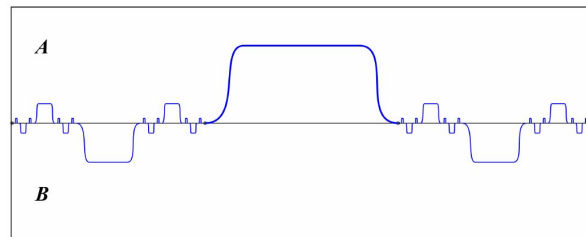
- Diffeological spaces, Frölicher spaces (exponential objects)
- Manifolds with (generalized) corners (flows on products and subsets of smooth manifolds)

2. Simplifications

- Not useful for us: gradient systems on surfaces, Morse systems, cobordisms, directed spaces
- Still useful for us: only trivially coupled systems, no zero behavior, no differential inclusions, restricting to nonblocking, deterministic systems



[De, Koditschek, 2018]



[Guralnik, 2020]

Concluding thoughts

1. Theorems of the form “* can be expressed as...” are not worth much in applications unless accompanied by meaningful machinery
2. Not every problem requires formalization/axiomatization
3. Working with practitioners continually is critical to making **useful** simplifying assumptions, abstractions, and interfaces

